Disciplined Approximate Computing: From Language to Hardware, and Beyond

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joint work with Adrian Sampson, Hadi Esmaelizadeh, Mike Ringenburg, Renee StAmant, Ben Ransford, Andre Baixo, Thierry Moreau, Dan Grossman, Mark Oskin (UW), Karin Strauss, Doug Burger, Todd Mytkowicz and Kathryn McKinley (Microsoft Research).



image, sound and video processing

image rendering

sensor data analysis, computer vision

simulations, games, search, machine learning



image, sound and video processing

# manna

# These applications consume a lot (most?)

# Often input data is inexact by nature (from sensors)

#### sensor data analysis, They have multiple acceptable outputs

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image, sound and video processing

# man

# These applications consume a lot (most?)

# Often input data is inexact by nature (from sensors)

#### sensor data analysis, They have multiple acceptable outputs

### They do not require "perfect execution"

simulations, games,

паснне

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# Notions of "approximation" have been around for a long time...

Floating point Lossy compression Iterative algorithms

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Resource usage (e.g., energy)

Sources of systematic accuracy loss:

Floating point

Lossy compression Iterative algorithms

Accuracy

- Unsound code transformations, ~2X
- Unreliable, probabilistic hardware (near/sub-threshold, etc.), ~5X
- Fundamentally different, inherently inaccurate execution models, "closer to physics" (e.g., neural networks, analog computing), ~10-100X







# But approximation needs to be done carefully... or...









#### Rapid

A fatal exception 8E has occurred at 8828:C8011E36 in UXD UHH(81) + 86818E36. The current application will be terminated. \* Press any key to terminate the current application. \* Press CTRL+ALT+DEL again to restart your computer. You will lose any unsaved information in all applications.

Press any key to continue \_

#### "Disciplined" approximate programming



### "Disciplined" approximate programming



Programmer has direct control of approximate/precise and the flow
 System is free to approximate as long as rules are obeyed

```
int p = 5;
@Approx int a = 7;
for (int x = 0..) {
    a += func(2);
    @Approx int z;
    z = p * 2;
    p += 4;
}
a /= 9;
p += 10;
socket.send(z);
write(file, z);
```

**Relaxed Algorithms** 

int p = 5; @Approx int a = 7; for (int x = 0..) { a += func(2); @Approx int z; z = p \* 2; p += 4; } a /= 9; p += 10; socket.send(z); write(file, z);

**Relaxed Algorithms** 

**Aggressive Compilation** 

```
int p = 5;
@Approx int a = 7;
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**Relaxed Algorithms** 

# Goal: support a wide range of approximation techniques with a **single unified** abstraction.

communication



/ariable-Accuracy ISA

Approximate Logic/Circuits



#### The plan for this talk

Application
Language
Compiler
Architecture
Circuits

### The plan for the rest of this talk





EnerJ: Approximate Data Types for Safe and General Low-Power Computation, PLDI 2011

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Separate critical and non-critical program components. *Analyzable statically.* 



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Separate critical and non-critical program components. *Analyzable statically.* 



### @Approx int a = ...; @Precise int p = ...;

EnerJ: Approximate Data Types for Safe and General Low-Power Computation, PLDI 2011

Separate critical and non-critical program components. *Analyzable statically.* 



## @Approx int a = ...; @Precise int p = ...; p = a; a = p;

EnerJ: Approximate Data Types for Safe and General Low-Power Computation, PLDI 2011

@Approx int a = ...; @Precise int p = ...;

- Operator overloading for approximate operations:
- Endorsement of approximate values:
- Dealing with implicit flows in control:
EnerJ: Approximate Data Types for Safe and General Low-Power Computation, PLDI 2011

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   operations: p + p; p + a; a + a;
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   **p** = **endorse(a);**
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- Operator overloading for approximate
   operations: p + p; p + a; a + a;
- Endorsement of approximate values:

   *p* = endorse(a);
- Dealing with implicit flows in control:





#### How good is my final output?

- Quality-of-Result (QoR)
- Application dependent
  - e.g, % of bad pixels, deviation from expected value, %
     of poorly classified images, car crashes, etc...



#### Specifying and checking QoR

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## res = computeSomething(); assert diff(res, res') < 0.1;</pre>

precise version of the result

#### Verifying quality expressions



Expressing and Verifying Probabilistic Assertions, PLDI'14

Can *react* – recompute or reduce approximation But needs to be cheap!

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Sampled precise re-execution

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Sampled precise re-execution Simple verification functions

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Sampled precise re-execution ·3> ● - ● **-** - <€? Simple verification functions

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Sampled Simple verification precise functions re-execution ·3> ● - ● 

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Sampled precise re-execution  $\bullet - \bullet < \epsilon?$  $\bullet - \bullet < \epsilon?$ 

Simple verification functions

Can react – recompute or reduce approximation But needs to be cheap!

Sampled Simple **Fuzzy Memoization** precise verification functions re-execution -**−**<ε?

 $< 2^{2}$ 

#### What about actual approximate execution?



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## Hardware support for disciplined approximate execution



Architecture Support for Disciplined Approximate Programming, ASPLOS 2012

#### @Approx float[] nums;



approximate data storage



approximate operations

### Relaxing the hardware-software interface



#### **Approximation-aware ISA**

ld	$0 \times 04$	r1	
ld	0x08	r2	
add	r1	r2	r3
st	0x0c	r3	

#### **Approximation-aware ISA**

 ld
 0x04 r1

 ld
 0x08 r2

 add.a
 r1
 r2
 r3

 st.a
 0x0c r3

#### **Approximation-aware ISA**

# ld 0x04 r1 ld 0x08 r2 add.a r1 r2 r3 st.a 0x0c r3

operations



storage

registers caches main memory

#### **Dual-voltage pipeline**

replicated functional units


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replicated functional units



#### 7-24% energy saved on average

(fft, game engines, raytracing, QR code readers, etc)

(scope: processor + memory)

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(fft, game engines, raytracing, QR code readers, etc)

(scope: processor + memory)

(though better implementations likely)

# Amdahl's law... damn!



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Benefit limited to what can be approximated
Instruction control can not be approximated

# How can we get rid of exact instruction bookkeeping?

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If behavior is approximate, why program it precisely?





#### Why Neural Networks as Approximate Accelerators?



Neural Acceleration of General-Purpose Approximate Programs, MICRO 2012 General-Purpose Code Acceleration with Limited-Precision Analog Computation, ISCA 2014

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**Find** an approximate program component



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**Find** an approximate program component

**Compile** the program and train a neural network





Program

**Find** an approximate program component

**Compile** the program and train a neural network

**Execute** on a fast Neural Processing Unit (NPU)

# **An example: Sobel filter**

}



edgeDetection()

ON ZO

@approx float grad(approx float[3][3] p) {

void edgeDetection(aImage &src, aImage &dst) { **for** (**int** y = ...) { for (int x = ...) { dst[x][y] =grad(window(src, x, y)); }

@approx float dst[][];

# An example: Sobel filter





```
@approx float grad(approx float[3][3] p) {
           Approximable
           Well-defined inputs and outputs
void edgeDetection(aImage &src,
                   aImage &dst) {
  for (int y = ...) {
    for (int x = ...) {
      dst[x][y] =
           grad(window(src, x, y));
```

@approx float dst[][];

#### **Empirically selecting** target code



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Each region of code leads to a different NN configuration.

# **Neural Processing Unit**



# **A digital NPU**



# **A digital NPU**



# How do the NNs look like in practice?



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# **Summary of results**

2.3x average speedup Ranges from 0.8x to 11.1x

**3.0x** average energy reduction for digital, ~**10x** for analog All benchmarks benefit

Quality loss below 10% in all cases Based on application-specific quality metrics

Just one possible design. Many others possible. Analog is where the big gains are likely (~10x+).

Key here is algorithmic transformation that enables new more efficient execution models.









#### Approximate mass storage with Flash and PCM

Approximate Storage in Solid State Memories [MICRO'13]

Cells wear out

over time

approximate otorage in oond otate memories [inforto Toj

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over time

Multi-level cells are slow or unreliable

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Use **worn-out** memory for **approximate** data instead of throwing it away.

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Cells **wear out** over time

Use **worn-out** memory for **approximate** data instead of throwing it away.

Multi-level cells are slow or unreliable

Trade off accuracy for performance/density in **multi-level cell** accesses.

## **Precise Multi-level Cells**



# **Approximate Multi-level Cells**


#### **Typical Trade-off in Multi-Level Cells**



#### Adding a New Trade-Off Axis















Configurable-quality wireless protocol. Quality automatically set by the data type.

# Neural Acceleration on a programmable SoC





### **Showing End-to-End benefit**

**Mobile Vision/Augmented Reality** Linux on Zynq SoC (ARM CPU + FPGA)



Applications can't take advantage of approximation opportunities

Programmers aren't able to write/debug/test approximate code

•Quality assurance problems

•Marketing reasons: "buy my flaky system!"

Applications can't take advantage of approximation opportunities

- Programmers aren't able to write/debug/test approximate
   Code

   [-] jtra 3 points 1 month ago
- •Quality assurance problems

- Good luck debugging that...
- •Marketing reasons: "buy my flaky system!"

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<sup>+</sup> Oh god, this sounds awful .. and it has enough potential to actually make things worse, if hardware vendors end up shoving it down our throats by force (i.e. people recalculating things over and over again, just to be safe; nonstandard/"unofficial" hardware that tries to work around the limitations in embedded devices, and other things like that).

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#### [-] MorePudding 16 points 1 month ago

There was a somewhat related post to this a few weeks ago here.

The basic idea was simple: if hardware suffers more transient failures as it gets smaller, why not allow software to detect erroneous computations and re-execute them? This idea seemed promising until John realized THAT IT WAS THE WORST IDEA EVER. Modern software barely works when the hardware is correct, so relying on software to correct hardware errors is like asking Godzilla to prevent Mega-Godzilla from terrorizing Japan.

### **Other ongoing effort**

- •Understanding specialization vs. approximation benefits
- Compiler-only approximation w/ unsound

transformations

- •HCI aspects: how do measure user satisfaction? do incentives matter in choosing quality?
- •Language support for QoR (quality of results,

probabilistic assertions)

- •**Tools** to help programmers w/ porting, testing and debugging
- •Exploring uses in **energy-harvesting**-based devices

#### <u>approxbench.org</u>

#### Conclusion

We need to exploit application properties and co-design hardwaresoftware for better efficiency.

Getting closer to physics might lead to very big efficiency gains.

Our goal is to exploit *approximate computing across the system*. (compute, storage, communication)

Key aspect is co-designing programming model with approximation techniques: *disciplined* approximate programming.

Early results encouraging. Approximate computing can potentially save our bacon in a post-Dennard era and be in the survival kit for dark silicon.

## Thanks!

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#### Microsoft<sup>®</sup> Research







